## THE CHARACTERIZATION OF EDGE-MAXIMAL CRITICALLY K-EDGE CONNECTED GRAPHS

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## ABSTRACT:

Let G be a simple graph on n vertices having edge-connectivity  $\kappa'(G) > 0$ . We say G is k-critical if  $\kappa'(G) = k$  and  $\kappa'(G - e) < k$  for every edge e of G. We denote by  $\mathcal{C}(n,k)$  the set of all k-critical graphs on n vertices. In this paper we prove that the maximum number of edges of a graph G in  $\mathcal{C}(n,k)$  to be: k(n-k) if  $n \ge 3k$ ; and  $\lfloor \frac{1}{8} (n+k)^2 \rfloor$ , if  $k + 1 \le n < 3k$ . Further, we characterise the extremal graphs in  $\mathcal{C}(n,k)$ .

#### 1. INTRODUCTION

All graphs considered in this paper are undirected, finite, loopless and have no multiple edges. For the most part our notation and terminology follows that of Bondy and Murty [2]. Thus G is a graph with vertex set V(G), edge set E(G),  $\nu$ (G) vertices,  $\varepsilon$ (G) edges, edge-connectivity  $\kappa'$ (G), and minimum degree  $\delta$ (G).  $K_n$  denotes the complete graph on n vertices,  $K_{m,n}$  the complete bipartite graph with bipartitioning sets of order m and n;  $C_n$  a cycle of length n. The join of disjoint graphs G and H, denoted by G v H, is the graph obtained by

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joining each vertex of G to each vertex of H. However, we denote the complement of G by  $\bar{G}$ .

We say a graph G is k-critical if  $\kappa'(G) = k$  and  $\kappa'(G-e) < k$  for every edge e of G. Observe that : every tree is 1-critical;  $C_n$  is 2-critical;  $K_1 \vee C_n$  is 3-critical;  $K_n$  is (n-1)-critical; and  $K_{m,n}$  is k-critical, where  $k = \min\{m,n\}$ . For fixed positive integers n and k, k-critical graphs on n vertices may not be unique. We denote by  $\mathcal{C}(n,k)$ the set of all k-critical graphs on n vertices. Let  $\mathcal{A}(n,k)$  denote the members of  $\mathcal{C}(n,k)$  that have every edge incident to at least one vertex of degree k.

Define  $\mathcal{B}(n,k) = \mathcal{C}(n,k) - \mathcal{A}(n,k)$ . We will establish that  $\mathcal{B}(n,k) = \phi$ for  $n \le 2k + 1$  and that  $\mathcal{B}(n,k) \ne \phi$  for  $n \ge 2(k+1)$ . An edge e = xy of G  $\in \mathcal{C}(n,k)$  will be called a **distinguished edge** if  $d_{G}(x) \ge k + 1$  and  $d_{G}(y) \ge k + 1$ . Thus every graph G in  $\mathcal{B}(n,k)$  contains at least one distinguished edge.

A graph  $G \in C(n,k)$  is called **edge-minimal (maximal)** if there is no other graph in C(n,k) having **less (more)** edges than G. We call G **r-semi-regular graph** if every vertex of G has degree r except one which has degree r + 1. We denote by H(n,t) a t-edge connected, t-regular (semi-regular) graph on n vertices for nt even (odd). Clearly this graph has  $\lceil \frac{nt}{2} \rceil$  edges.

In [3] we proved that if  $G \in C(n,k)$  then  $\delta(G) = k$ . We also proved that  $G \in C(n,k)$  if and only if there are exactly k edge-disjoint paths joining any two adjacent vertices of G. So it is obvious that for  $k \neq 1$ , a graph  $G \in C(n,k)$  is edge-minimal if and only if G =H(n,k). The edge-maximal members of C(n,k) are not as easily described. Indeed, their structure is much more complex. In [3] we considered the problem of determining the maximum number of edges for a graph  $G \in \mathcal{A}(n,k)$ . We proved that this number is equal to : k(n-k) for  $n \ge 3k$ ; and  $\left\lfloor \frac{1}{8} (n+k)^2 \right\rfloor$ , for  $k + 1 \le n < 3k$ . In this paper we will prove that this result is true for an edge-maximal graph in  $\mathcal{C}(n,k)$ . We also prove that there is no distinguished edge in an edge- maximal graph of  $\mathcal{C}(n,k)$  for  $k \ne 1$ . The edge-maximal graphs in  $\mathcal{C}(n,k)$  are completely characterised.

### 2. FUNDAMENTAL LEMMAS

We define an edge-cut of a graph G as a subset of E(G) of the form  $(V_1, \overline{V}_1)$ , where  $V_1$  is a nonempty proper subset of V(G). A k-edge cut is an edge cut of k elements.

Suppose G is a k-edge connected graph having two k-edge-cut sets, say  $E_1$  and  $E_2$ . Removing  $E_1 \cup E_2$  from G yields a graph G' having three or four components. We will show that if G' has four components then every component is separated by a k-edge cut set from G; but if it has three components then at least two of them are separated by k-edge cut sets from G.

Lemma 2.1 : Let G be a k-edge connected graph on n vertices with k-edge cuts  $E_1$  and  $E_2$ . If G -  $(E_1 \cup E_2)$  consists of four components,  $G_1$ ,  $G_2$ ,  $G_3$  and  $G_4$ , then for every i,  $1 \le i \le 4$ ,  $(V_1, \overline{V}_1)$  is a k-edge cut in G, where  $V_i = V(G_i)$ .

Proof : Without loss of generality we assume that

$$\mathbf{E}_{1} = (\mathbf{V}_{1} \cup \mathbf{V}_{2}, \mathbf{V}_{3} \cup \mathbf{V}_{4})$$

and

$$\mathbf{E}_2 = (\mathbf{V}_1 \cup \mathbf{V}_4, \mathbf{V}_2 \cup \mathbf{V}_3).$$

We let  $e_{ij}$  denote the number of edges in G between  $G_i$  and  $G_j$  (see Figure 2.1).

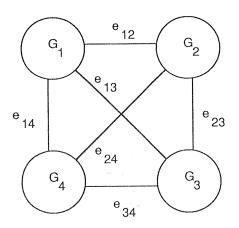


Figure 2.1

Observe that

$$k = |E_1| = e_{13} + e_{14} + e_{23} + e_{24}$$
(1)

and

$$k = |E_2| = e_{12} + e_{13} + e_{24} + e_{34}$$
 (2)

Further, since  $\kappa'(G) = k$  we must have

$$e_{12} + e_{13} + e_{14} \ge k$$
, (3)

$$e_{12} + e_{23} + e_{24} \ge k$$
, (4)

$$e_{13} + e_{23} + e_{34} \ge k$$
, (5)

and

$$e_{14} + e_{24} + e_{34} \ge k$$
 (6)

Now (1) and (3), (2) and (4), (1) and (5), and (2) and (6) respectively imply :

$$e_{12} \ge e_{23} + e_{24}$$
, (7)

$$e_{23} \ge e_{13} + e_{34}$$
, (8)

$$e_{34} \ge e_{14} + e_{24}$$
, (9)

and 
$$e_{14} \ge e_{12} + e_{13}$$
 (10)

Since  $e_{ij} \ge 0$ , inequalities (7) to (10) together imply that  $e_{12} = e_{23} = e_{34} = e_{14}$  and hence  $e_{13} = e_{24} = 0$ . Now equations (1) and (2) give  $e_{14} = e_{23} = e_{12} = e_{34} = \frac{1}{2} k$ .

This completes the proof of Lemma 2.1.

Note that the above proof yields the following result.

**Corollary 2.1**: Let G be a k-edge connected graph with k-edge cuts  $E_1$  and  $E_2$ . If G -  $(E_1 \cup E_2)$  consists of four components, then k is even and  $E_1 \cap E_2 = \phi$ . Further, if k is odd G -  $(E_1 \cup E_2)$  consists of three components.

Our next lemma considers the case when G -  $E_1^{}$   $\cup$   $E_2^{}$  has 3 components and k is odd or even.

**Lemma 2.2**: Let G be a k-edge-connected graph on n vertices with k-edge cuts  $E_1$  and  $E_2$  having t edges in common. If  $G - (E_1 \cup E_2)$  consists of 3 components  $G_1, G_2$  and  $G_3$ , then  $t \leq \frac{1}{2}$  k with equality holding only if  $(V_1, \overline{V}_1)$  is a k-edge-cut set in G for each i,  $1 \leq i \leq 3$ , where  $V_1 = V(G_1)$ . Furthermore, if  $t < \frac{1}{2}k$  then  $(V_1, \overline{V}_1)$  is a k-edge cut set in G for exactly two of the i's.

Proof : Without loss of generality we assume that

$$E_1 = (V_1, V_2 \cup V_3)$$
 and  $E_2 = (V_1 \cup V_2, V_3)$ .

As in the above proof we let  $e_{ij}$  denote the number of edges in G between  $G_i$  and  $G_j$  (see Figure 2.2). Note that  $t = e_{13}$ .

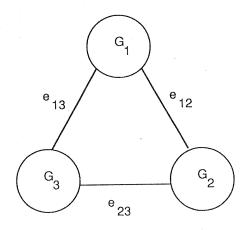


Figure 2.2

We have  $e_{12} + e_{13} = k$  and  $e_{13} + e_{23} = k$ , hence  $e_{12} = e_{23}$ . Now since  $\kappa'(G) = k$  we have

and thus  $e_{12} = e_{23} \ge \frac{1}{2} k$ . This implies that  $e_{13} \le \frac{1}{2} k$ . Furthermore, if  $e_{13} = \frac{1}{2} k$ , then

$$e_{12} = e_{13} = e_{23} = \frac{1}{2}k$$

Observe that if  $e_{13} < \frac{1}{2}k$ , then  $e_{12} + e_{23} > k$ . This completes the proof of the lemma.

Our next lemma is essential for proving our main result.

Lemma 2.3 : Let G be a graph on n vertices and let k be a positive integer less than n. If every edge of G is incident to at least one vertex of degree k or less, then G has at most  $\varepsilon(n,k)$  edges, where

$$\varepsilon(n,k) = \begin{cases} k(n-k) , & \text{if } n \ge 3k \\ \\ \\ \lfloor \frac{1}{8} (n+k)^2 \rfloor , & \text{if } k+1 \le n < 3k \end{cases}$$
(11)

**Proof**: Let G be a graph on n vertices with every edge incident to at least one vertex of degree k or less. We denote by  $A = \{v \in V(G) : d_{\overline{G}}(v) \le k\}$  and  $\overline{A} = V(G) - A$ . Let  $|A| = n_1$  and  $|\overline{A}| = n_2 = n - n_1$ . Since every vertex x of G in  $\overline{A}$ , has degree at least k+1, and  $E(G[\overline{A}]) = \phi$ , we have  $n_1 \ge k+1$ .

If  $n_1 \le n-k$ , then  $n_2 = n-n_1 \ge k$ , so the maximum number of edges of G is obtained when there is no edge e = xy of G such that  $x, y \in A$ . Hence  $\epsilon(G) \le n_1 k \le (n-k)k$ .

If  $n_1 \ge n-k$ , then  $n_2 \le k$ . Simple counting gives  $\varepsilon(G) \le n_1(n-n_1) + \lfloor \frac{1}{2} n_1(k-n+n_1) \rfloor$   $= \lfloor n_1(n-n_1) + \frac{1}{2} n_1(k-n+n_1) \rfloor$  $= \lfloor \hat{g}(n_1) \rfloor = g(n_1)$ .

For  $n \ge 3k$ ,  $\hat{g}(n_1)$  is a decreasing function of  $n_1$ , and so

$$\max_{n_1} \{ \hat{g}(n_1) \} = \lfloor \hat{g}(n-k) \rfloor$$
$$= (n-k)k$$

Since  $\hat{g}(n_1)$  is monotonically increasing when  $n_1 \leq \lfloor \frac{n+k}{2} \rfloor$  and monotonically decreasing when  $n_1 \geq \lceil \frac{n+k}{2} \rceil$ , we have for n < 3k,

$$\underset{n_{1}}{\overset{\text{max}}{\text{max}}} \{ G(n_{1}) \} = \max \{ g(\lfloor \frac{n+k}{2} \rfloor), g(\lceil \frac{n+k}{2} \rceil) \}$$
$$= \lfloor \frac{1}{8} (n+k)^{2} \rfloor .$$

Now since

$$(n+k)^2 - 8k(n-k) = (n-3k)^2 \ge 0$$

and k(n-k) is integer, we have

$$\left\lfloor \frac{1}{8} (n+k)^2 \right\rfloor \ge k(n-k) .$$

Therefore

This completes the proof of Lemma 2.3.

The following result was proved in [3]. It is actually a special case of Lemma 2.3.

**Corollary 2.2 :** If  $G \in \mathcal{A}(n,k)$ , then the maximum number edges of G is  $\varepsilon(n,k)$ , where  $\varepsilon(n,k)$  is given by (11).

Our next result considers the class  $\mathcal{B}(n,k)$ .

Lemma 2.4 : Let  $G \in \mathcal{B}(n,k) = \mathcal{C}(n,k) - \mathcal{A}(n,k)$ . Then

$$\nu(G) \geq \begin{cases} 7 & , \text{ if } k = 2 \\ \\ 2(k+1) & , \text{ otherwise } \end{cases}$$

Furthermore, this bound is sharp.

**Proof** : By definition, the graph  $G \in \mathcal{B}(n,k)$  contains a distinguished edge  $e_1$ , say. Now let  $E_1$  be a k-edge cut of G containing  $e_1$ . Let  $G_1$ and  $G_2$  be the components of  $G - E_1$  and suppose that  $|V(G_1)| = n_1$ , i = 1,2. Without loss of generality suppose  $n_1 \leq n_2$ . We will show that  $n_1 \geq k+1$ . In fact, since  $e_1$  is a distinguished edge of G joining  $G_1$  and  $G_2$ , there is a vertex, x say, of  $G_1$  in G with  $d_G(x) \geq k+1$ . Since G is k-critical then  $\delta(G) = k$ . Hence

$$\sum_{\mathbf{v}\in \mathbf{G}_{1}} \mathbf{d}_{\mathbf{G}}(\mathbf{v}) \geq \mathbf{n}_{1}\mathbf{k}+1 \ .$$

186

Now if  $n_1 \leq k$ , we have

$$\sum_{\mathbf{v}\in\mathbf{G}_{1}} \mathbf{d}_{\mathbf{G}}(\mathbf{v}) = \sum_{\mathbf{v}\in\mathbf{G}_{1}} \mathbf{d}_{\mathbf{G}_{1}}(\mathbf{v}) + \mathbf{k}$$
$$\leq \mathbf{n}_{1}(\mathbf{n}_{1}-1) + \mathbf{k}$$
$$\leq \mathbf{k}(\mathbf{n}_{1}-1) + \mathbf{k} = \mathbf{n}_{1}\mathbf{k}$$

contradicting the above fact. Hence  $n_1 \ge k+1$ . Now since  $n_2 \ge n_1$ , we must have

$$n = n_1 + n_2 \ge 2(k+1)$$
.

For k = 2, a straight forward case analysis establishes that, the graph G formed from  $K_3$  and  $C_4$  by adding two edges joining a vertex of  $K_3$  and two nonadjacent vertices of  $C_4$ , is the member of  $\mathcal{B}(n,k)$  with the smallest number of vertices. This establishes the lower bound on  $\nu(G)$ . For  $k \neq 2$  we establish the sharpness of this bound by construction.

For k odd, we construct the graph  $G \in \mathcal{B}(2k+2,k)$  as follows. Take  $G_1 = K_1 \vee H(k,k-2)$ . Since k(k-2) is odd, there is a vertex, x say, of H(k,k-2) with  $d_H(x) = k-1$ . We form G by taking two copies  $G'_1$  and  $G''_1$  of  $G_1$  and adding a perfect matching between the vertices of H' and H" with x'x" an edge in this matching. Observe that G is a k-critical graph, on 2(k+1) vertices, and  $d_G(x') = d_G(x'') = k+1$ . So x'x" is a distinguished edge of G. Thus  $G \in \mathcal{B}(n = 2k+2,k)$ .

For k even,  $k \neq 2$ , we construct the graph  $G \in \mathcal{B}(2k+2,k)$  as follows. Take  $G_1 = \overline{K}_2 \vee K_{k-1}$ , and denote the vertices of  $\overline{K}_2$  by x and y. Let  $G_2 = K_2 \vee H(k-1,k-3)$ . Since (k-1)(k-3) is odd there is a vertex, z say, of H(k-1,k-3) with  $d_H(z) = k-2$ . Form the graph G from  $G_1$  and  $G_2$  by joining z to x and y, and joining all other vertices of H to exactly

187

one of x or y so that x and y have the same degree in G. Observe that G is a k-critical graph on 2k+2 vertices, with  $d_{G}(x) = d_{G}(y) = \left\lceil \frac{k-1}{2} \right\rceil + k-1 \ge k+1$  and  $d_{G}(z) = k+2$ . Therefore  $G \in \mathcal{B}(n = 2k+2, k)$ . This completes the proof of the lemma.

Lemma 2.5 : For positive integers n and k with k > 1 let

$$a(n,k) = \begin{cases} \frac{1}{2} n(n-1) , & \text{if } 1 \le n \le k \\\\ \varepsilon(n,k) , & \text{otherwise.} \end{cases}$$

where  $\varepsilon(n,k)$  is given by (11). If  $n_1 \ge k+1$ , then

$$a(n_1,k) + a(n_2,k) + k \le a(n_1 + n_2, k)$$
 (12)

unless k+1  $\leq n_1 < 3k-1$  and  $n_2 = 1$  in which case  $a(n_1,k) + a(n_2,k) + \frac{1}{2}k \leq a(n_1+n_2,k) .$ 

**Proof** : We define the function

 $f(n_1, n_2, k) = a(n_1, k) + a(n_2, k) + k.$ 

We consider two cases according to the value of n.

Case (a): 
$$n_1 \ge 3k$$

If  $n_2 \ge 3k$ , then

$$f(n_1, n_2, k) = k(n_1 - k) + k(n_2 - k) + k$$
  
=  $(n_1 + n_2 - k)k - (k^2 - k)$   
<  $(n_1 + n_2 - k)k = a(n_1 + n_2, k)$ 

This proves the lemma for the case  $n_2 \ge 3k$ . If  $k+1 \le n_2 < 3k$ , then

$$f(n_1, n_2, k) = k(n_1 - k) + \left\lfloor \frac{1}{8}(n_2 + k)^2 \right\rfloor + k$$
$$\leq k(n_1 - k) + \frac{1}{8}(n_2 + k)^2 + k$$

$$= a(n_1 + n_2, k) + \frac{1}{8}(n_2 + k)^2 + k - n_2 k$$

Now it is a simple algebraic exercise to show that

$$\frac{1}{8} (n_2 + k)^2 + k - n_2 k < 0 .$$

Hence

$$f(n_1, n_2, k) < a(n_1 + n_2, k) \text{ for } k + 1 \le n_2 < 3k$$

The only remaining case is  $n_2 \leq k$ .

For  $n_2 \leq k$ , we have

$$f(n_1, n_2, k) = k(n_1 - k) + \frac{1}{2}n_2(n_2 - 1) + k$$
$$= a(n_1 + n_2, k) + \frac{1}{2}n_2(n_2 - 1) + k - kn_2$$

Now for  $n_2 \leq k$  the function

$$h(n_2, k) = \frac{1}{2} n_2(n_2 - 1) + k - kn_2$$

is monotonically decreasing in n<sub>2</sub>. Hence

$$h(n_{1},k) \leq h(1,k) = 0$$
.

Therefore

$$f(n_1, n_2, k) \le a(n_1 + n_2, k)$$

with equality possible only if  $n_2 = 1$ .

**Case (b):**  $k + 1 \le n_1 < 3k$ .

We may assume that  $n_2 < 3k$  as otherwise we can, by interchanging  $n_1$  and  $n_2$ , apply the above argument.

If  $n_2 \ge k + 1$ , then

$$f(n_1, n_2, k) = \left\lfloor \frac{1}{8} (n_1 + k)^2 \right\rfloor + \left\lfloor \frac{1}{8} (n_2 + k)^2 \right\rfloor + k$$
$$\leq \left\lfloor \frac{1}{8} \left\{ (n_1 + k)^2 + (n_2 + k)^2 \right\} + k \right\rfloor$$

We first consider the case when  $n_1 + n_2 \ge 3k$ . In this case  $a(n_1 + n_2, k) = k(n_1 + n_2 - k)$ . Now the function

$$h(n_1, n_2, k) = \frac{1}{8} \{(n_1 + k)^2 + (n_2 + k)^2\} + k - k(n_1 + n_2 - k)$$

is for k + 1  $\leq$  n<sub>2</sub> < 3k, monotonically decreasing in n<sub>2</sub> so it is maximum when n<sub>2</sub> is as small as possible. Hence, since n<sub>2</sub>  $\geq$  k + 1

$$\begin{array}{l} \max \\ n_2 \\ n_2 \\ \end{array} \left\{ h(n_1, n_2, k) \right\} = h(n_1, \ k + 1, \ k) \\ \\ = \frac{1}{8} \left\{ (n_1 + k)^2 + (2k + 1)^2 \right\} + k - k(n_1 + 1) \\ \\ \\ = \frac{1}{8} (n_1 - 3k)^2 - (2k - 1)^2 + 2 \right\} .$$

Since for  $k + 1 \le n_1 < 3k$ , we have  $(n_1 - 3k)^2 \le (1 - 2k)^2$  we conclude that

$$\max_{\substack{n_1, n_2}} \{h(n_1, n_2, k)\} = \frac{1}{8} \{(1 - 2k)^2 - (2k - 1)^2 + 2\}$$
$$\leq \frac{1}{4}.$$

Thus we have

$$\frac{1}{8} \{ (n_1 + k)^2 + (n_2 + k)^2 \} + k - k(n_1 + n_2 - k) \le \frac{1}{4} .$$

Now since  $a(n_1 + n_2, k) = k(n_1 + n_2 - k)$  is an integer, we have

$$\left\lfloor \frac{1}{8} \left\{ \left( n_{1} + k \right)^{2} + \left( n_{2} + k \right)^{2} \right\} + k \right\rfloor - k \left( n_{1} + n_{2} - k \right) \le 0.$$

Hence

$$f(n_1, n_2, k) \le a(n_1 + n_2, k)$$
.

This completes the proof for the case  $n_1 + n_2 \ge 3k$ .

We now consider the case  $n_1 + n_2 < 3k$ . For this case we have

$$a(n_1 + n_2, k) = \left\lfloor \frac{1}{8}(n_1 + n_2 + k)^2 \right\rfloor.$$

The function

$$h(n_1, n_2, k) = \frac{1}{8} \{(n_1 + k)^2 + (n_2 + k)^2\} + k - \frac{1}{8} (n_1 + n_2 + k)^2$$
$$= \frac{1}{8} \{k^2 + 8k - 2n_1n_2\}$$
$$\leq \frac{1}{8} \{k^2 + 8k - 2(k + 1)^2\}$$
$$= \frac{1}{8} \{-k^2 + 4k - 2\}$$
$$\leq \frac{1}{8} (as \ k \ge 3).$$

(Note that this case does not happen for k = 2). Again as  $f(n_1, n_2, k)$ and  $a(n_1 + n_2, k)$  are integers

$$f(n_1, n_2, k) - a(n_1 + n_2, k) \le 0$$

as required.

The only remaining case is  $n_2 \leq k$ . We have

$$f(n_1, n_2, k) = \left\lfloor \frac{1}{8}(n_1 + k)^2 \right\rfloor + \frac{1}{2} n_2(n_2 - 1) + k$$
$$\leq \frac{1}{8}(n_1 + k)^2 + \frac{1}{2} n_2(n_2 - 1) + k.$$

Now

We first consider the case  $n_1 + n_2 \ge 3k$ . Then the function

$$h(n_1, n_2, k) = \frac{1}{8}(n_1 + k)^2 + \frac{1}{2}n_2(n_2 - 1) + k - k(n_1 + n_2 - k)$$

is, for  $1 \le n_2 \le k$ , monotonically decreasing in  $n_2$ . Thus, since  $n_2 \ge 3k - n_1 \ge 1$ ,

$$\begin{array}{l} \max_{n_2} \{h(n_1, n_2, k)\} = h(n_1, 3k - n_1, k) \\ \\ = \frac{1}{8} \{5n_1^2 - (22k - 4)n_1 + 17k^2 - 4k\} \\ \\ = \hat{h}(n_1, k) , \end{array}$$

and

$$\frac{\max_{n} \{\hat{h}(n_{1}k)\} = \max \{\hat{h}(k+1, k), \hat{h}(3k-1, k)\}$$

$$= \hat{h}(k+1, k)$$

$$= \frac{1}{8}(-12k + 9) < -1\frac{7}{8}$$

Hence

$$f(n_1, n_2, k) < a(n_1 + n_2, k)$$
.

Now we consider the case  $n_1 + n_2 < 3k$ .

The function

$$h(n_1, n_2, k) = \frac{1}{8}(n_1 + k)^2 + \frac{1}{2}n_2(n_2 - 1) + k - \frac{1}{8}(n_1 + n_2 + k)^2$$
$$= \frac{3}{8}n_2^2 - \frac{1}{4}(n_1 + k + 2)n_2 + k$$

is, for  $k + 1 \le n_1 < 3k$ , monotonically decreasing in  $n_1$ . Hence

$$\max_{n_{1}} \{h(n_{1}, n_{2}, k)\} = h(k+1, n_{2}, k)$$
$$= \frac{3}{8}n_{2}^{2} - \frac{1}{4}(2k+3)n_{2} + k$$
$$\leq 0 \text{ for } n_{2} \geq 2.$$

Thus for  $n_2 \ge 2$  we have proved the lemma.

The only remaining case is  $n_2 = 1$  and  $k+1 \le n_1 < 3k-1$ . For this case we have

 $a(n_1 + n_2, k) = \left\lfloor \frac{1}{8}(n_1 + 1 + k)^2 \right\rfloor$ 

and

$$f(n_1, n_2, k) = \left\lfloor \frac{1}{8}(n_1 + k)^2 \right\rfloor + k$$
.

Now consider

$$f^{*}(n_{1}, n_{2}, k) = f(n_{1}, n_{2}, k) - \frac{1}{2}k$$
$$= \left\lfloor \frac{1}{8}(n_{1} + k)^{2} \right\rfloor + \frac{1}{2}k$$
$$\leq \frac{1}{8}(n_{1} + k)^{2} + \frac{1}{2}k.$$

We have

$$h(n_1, n_2, k) = \frac{1}{8}(n_1 + k)^2 + \frac{1}{2}k - \frac{1}{8}(n_1 + 1 + k)^2$$
$$= \frac{1}{2}k - \frac{1}{4}(n_1 + k) + \frac{1}{8}$$
$$\le \frac{1}{2}k - \frac{1}{4}(2k + 1) + \frac{1}{8} < 0$$

Hence

$$f^*(n_1, n_2, k) \le a(n_1 + n_2, k)$$
 as required.

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This completes the proof of the lemma.

**Lemma 2.6**: For  $3 \le k+1 \le n_i < 3k$ , i = 1, 2

$$\varepsilon(n_1,k) + \varepsilon(n_2,k) + 2k \le \varepsilon(n_1 + n_2 + 1,k)$$
<sup>(13)</sup>

where  $\varepsilon(n,k)$  is given by (11).

Proof : Let

$$h(n_{1}, n_{2}, k) = \varepsilon(n_{1}, k) + \varepsilon(n_{2}, k) + 2k - \varepsilon(n_{1} + n_{2} + 1, k).$$

We distinguish two cases according to the value of  $n_1 + n_2 + 1$ .

**Case (a):**  $n_1 + n_2 + 1 \ge 3k$ .

We have

$$h(n_1, n_2, k) = \left\lfloor \frac{1}{8} (n_1 + k)^2 \right\rfloor + \left\lfloor \frac{1}{8} (n_2 + k)^2 \right\rfloor + 2k$$
  
-  $(n_1 + n_2 + 1 - k)k$   
$$\leq \frac{1}{8} \{ (n_1 + k)^2 + (n_2 + k)^2 \} + 2k - k(n_1 + n_2 + 1 - k)$$
  
=  $\hat{h} (n_1, n_2, k).$ 

The function  $\hat{h}(n_1, n_2, k)$  is, for  $k+1 \le n_1 < 3k$ , monotonically decreasing in  $n_1$ . So it is maximum when  $n_1$  is as small as possible. Hence, since  $n_1 \ge k+1$ 

$$\max_{n_{1}} \{\hat{h}(n_{1}, n_{2}, k)\} = \hat{h}(k+1, n_{2}, k)$$

$$= \frac{1}{8} \{(2k + 1)^{2} + (n_{2} + k)^{2}\}$$

$$+ 2k - k(n_{2} + 2)$$

$$= \frac{1}{8} \{(n_{2} - 3k)^{2} - (2k - 1)^{2} + 2\}.$$

Now for  $k + 1 \le n_2 < 3k$ , we can show that

$$(n_2 - 3k)^2 \le (1 - 2k)^2$$

Hence we have

$$\max_{\substack{n_1, n_2 \\ n_1, n_2}} \{\hat{h}(n_1, n_2, k) = \frac{1}{8} \{(1-2k)^2 - (2k-1)^2 + 2\} = \frac{2}{8}.$$

Since  $h(n_1, n_2, k)$  is an integer, we conclude that  $h(n_1, n_2, k) \le 0$ . This proves the lemma for this case.

**Case (b)**:  $n_1 + n_2 + 1 < 3k$ .

Note that this case occurs only when k > 3. Here we have

$$h(n_{1}, n_{2}, k) = \left\lfloor \frac{1}{8}(n_{1} + k)^{2} \right\rfloor + \left\lfloor \frac{1}{8}(n_{2} + k)^{2} \right\rfloor + 2k - \left\lfloor \frac{1}{8}(n_{1} + n_{2} + 1 + k)^{2} \right\rfloor$$

Consider the function

$$\hat{h}(n_1, n_2, k) = \frac{1}{8} \{ (n_1 + k)^2 + (n_2 + k)^2 - (n_1 + n_2 + 1 + k)^2 \} + 2k$$
$$= \frac{1}{8} \{ -2 (n_1 + n_2) - 2n_1n_2 + k^2 + 14k - 1 \}.$$

Since  $\hat{h}(n_1, n_2, k)$  is a monotonically decreasing function in  $n_2$ , and  $n_2 \ge k+1$ , we have

$$\max_{n_{2}} \{\hat{h}(n_{1}, n_{2}, k)\} = \hat{h}(n_{1}, k+1, k)$$

$$= \frac{1}{8}\{k^{2} + 12k - 3 - 2(2+k)n_{1}\}$$

$$\leq \frac{1}{8}\{ -k^{2} + 6k - 7\} \text{ (since } n_{1} < 3k\}$$

$$\leq \frac{1}{8} \text{ (as } (k \geq 4).$$

Consequently  $h(n_1, n_2, k) \le 0$ , as required. This completes the proof of Lemma 2.6.

# 3. MAIN RESULTS

The following terminology is useful in the proof of our main theorem. Let  $G \in \mathcal{C}(n,k)$ . For a subset U of V(G) we write  $\varepsilon(U,\overline{U})$  for the number of edges between U and  $\overline{U}$ .

When  $\varepsilon(U,\overline{U}) = k$  we call U a **segment** of G (note that  $\overline{U}$  is also a segment). A consequence of Lemmas 2.1 and 2.2 is the following result.

**Lemma 3.1:** If A and B are two segments such that  $B \cap \overline{A} \neq \phi$ , then either  $A \subseteq B$  or  $B \cap \overline{A}$  is a segment.

For  $X \subseteq V(G)$  we denote the subgraph of G induced by the vertices in X by G[X].

Theorem 3.1 : Let G be an edge maximal graph in C(n,k). Then

$$\varepsilon(G) = \begin{cases} k(n-k), & \text{if } n \ge 3k \\ \\ \\ \lfloor \frac{1}{8}(n+k)^2 \rfloor, & \text{if } k+1 \le n < 3k \end{cases}.$$

**Proof**: If G has no distinguished edge, then the result coincides with Corollary 2.2 and we have nothing to prove. So suppose G contains at least one distinguished edge.

Choose a k-edge cut  $E_1$  containing a distinguished edge,  $e_1$  say, such that  $G - E_1$  contains a component,  $G_1$  say, having no distinguished edge. That such an  $E_1$  and  $e_1$  exists follows from Lemmas 2.1 and 2.2. Let  $G_2$  be the other component of  $G - E_1$  and let  $n_1 = |V(G_1)|$ , i = 1, 2. As in the proof of Lemma 2.4, since  $G_1$  and  $G_2$  each contain a vertex of degree k+1 we have  $n_1 \ge k+1$  and  $n_2 \ge k+1$ . Since  $G_1$  has no distinguished edge of G, we have by Lemma 2.3,  $\varepsilon(G_1) \le \varepsilon(n_1, k)$ . Now if  $\varepsilon(G_2) \le \varepsilon(n_2, k)$ , then

$$\varepsilon(G) = \varepsilon(G_1) + \varepsilon(G_2) + k$$
  

$$\leq \varepsilon(n_1, k) + \varepsilon(n_2, k) + k$$
  

$$\leq \varepsilon(n_1 + n_2, k) \qquad (Lemma 2.5)$$
  

$$= \varepsilon(n, k) ,$$

as required. Thus we may assume that  $\varepsilon(G_2) > \varepsilon(n_2,k)$ . Then, by Lemma 2.3, G<sub>2</sub> contains at least one distinguished edge.

Our strategy is to partition the vertices of  $G_2$  into sets and then

apply Lemmas 2.5 and 2.6 to count the edges of G. Observe that if e' is a distinguished edge of G in  $G_2$  and E' is a k-edge cut of G containing e', then by Lemma 3.1 there exists a segment S' such that S'  $\cap V(G_1) = \phi$ . Further, Lemmas 2.1 and 2.2 ensure that we can choose e' and E' such that G[S'] contains no distinguished edge of G. Note that  $|S'| \ge k+1$  and by Lemma 2.3  $\varepsilon(G[S']) \le \varepsilon(|S'|, k)$ .

Let T denote the largest segment of G such that  $T \cap V(G_1) = \phi$  and  $\varepsilon(G[T]) \leq \varepsilon(|T|,k)$ . That such a T exists follows from the existence of S'. Since  $\varepsilon(G_2) > \varepsilon(n_2,k)$ ,  $T \neq V(G_2)$ . Let T' =  $V(G_2) \setminus T$ . If  $\varepsilon(G[T'] \leq \alpha(|T'|,k)$ , then

$$\varepsilon(G) \le \varepsilon(G_{1}) + \varepsilon(G[T]) + \varepsilon(G[T']) + 2k$$
  
$$\le a(n_{1},k) + a(|T|,k) + a(|T'|,k) + 2k$$
  
$$\le a(n_{1} + |T| + |T'|,k) \quad (Lemmas 2.5 \& 2.6)$$
  
$$= \varepsilon(n,k),$$

as required. Hence we may assume that  $\varepsilon(G[T'],k) > a(|T'|,k)$ . The definition of a(n,k) then implies that  $|T'| \ge k+1$ . We now partition the set T'.

We can find a distinguished edge  $e_1$  of G in G[T'] and a segment  $S_1$ of G such that  $e_1$  is in the cut  $(S_1, \overline{S}_1)$ ,  $S_1 \cap V(G_1) = \phi$  (so  $S_1 \cap \overline{T} = S_1$  $\cap T'$ ) and G  $[S_1 \cap \overline{T}]$  does not contain a distinguished edge. That this can be done follows from lemmas 2.1, 2.2 and 3.1. Continuing in this way we partition T' into sets  $B_1, B_2, \dots, B_t$  such that

$$B_{1} = S_{1} \wedge \overline{T} , \quad B_{1} = (S_{1} \setminus \bigcup S_{j}) \wedge \overline{T}$$

$$j=1$$

and each subgraph  $G[B_i]$  contains no distinguished edge of  $G_1$ .

197

Hence, by Lemma 2.3,  $\varepsilon(G[B_i] \le \varepsilon(|B_i|,k) \text{ when } |B_i| \ge k+1$ . Obviously  $\varepsilon(G[B_i]) \le \frac{1}{2}|B_i|(|B_i| -1)$ . Hence  $\varepsilon(G[B_i]) \le \alpha(|B_i|,k)$  for every i. If  $|B_i| \ge k+1$  for some i, with no loss of generality say i=1, then since  $|V(G_i)| + |T| + |B_i| \ge 3(k+1)$ , we have by Lemma 2.5

$$\varepsilon(G) \leq \varepsilon(G_1) + \varepsilon(G[T]) + \varepsilon(G[B_1] + 2k + \sum_{i=2}^{t} \varepsilon(G[B_i]) + (t-1)k$$

$$\leq a(|V(G_1)| + |T| + |B_1|,k) + \sum_{i=2}^{t} a(|B_i|,k) + (t-1)k$$

 $\leq a(|V(G_1)| + |T| + |B_1| + \dots + |B_t|,k)$ 

 $= \epsilon(n,k),$ 

as required. Thus assume that  $|B_1| \le k$  for every i. Then  $B_1$  is not a segment. Now by Lemma 3.1 T  $\le S_1$  and  $S_1 = T \cup B_1$ .

Suppose that  $|B_i| \ge 2$  for some i. With no loss of generality let  $|B_i| \ge 2$ . Then

$$\varepsilon(G[T \cup B_1]) \leq \varepsilon(G[T],k) + \varepsilon (G[B_1],k) + k$$

 $\leq a(|T|,k) + a(|B_1|,k) + k$ 

 $\leq a(|T \cup B_1|,k)$  (Lemma 2.5)

contradicting the choice of T as  $T \cup B_1$  is a segment of G. Thus  $|B_1| = 1$  for each i and further no  $B_1$  is a segment. Hence, since  $|T'| \ge k+1$ ,  $t \ge 2$ . By Lemma 3.1,  $S_1 \cup B_2$  is a segment of G. Now since  $|B_1 \cup B_2| = 2$  we have

$$\varepsilon(G[S_1 \cup B_2]) = \varepsilon(G[T \cup B_1 \cup B_2])$$

$$\leq \varepsilon(G[T]) + \varepsilon(G[B_1 \cup B_2]) + k$$

$$\leq a(|T|,k) + a(|B_1 \cup B_2|,k) + k$$

$$\leq a(|T| + |B_1 \cup B_2|,k) , \qquad (Lemma 2.5)$$

$$= a(|T \cup B_1 \cup B_2|,k) = a(|S_1 \cup B_2|,k),$$

contradicting the choice of T. This completes the proof of the theorem.  $\hfill \Box$ 

In the following result we will prove that an edge maximal graph in  $\mathcal{C}(n, k \neq 1)$  can not have a distinguished edge. We will make use of the following remarks in the proof of our result.

**Remark 3.1 :** It can be shown that equality in (12) holds only if one of the following conditions is satisfied :

(i)	$n_1 \ge 3k-1$	and	$n_2 = 1$ ; for every k
(ii)	$n_1 = n_2 = k+1$	and	$2 \leq k \leq 4$
(iii)	$n_1 = k+1$	and	$n_2 = 2$ ; for every k
(iv)	$n_1 = 4$	and	$n_2 = k = 3$
(v)	$n_1 = k+2$	and	$n_2 = 2; k \text{ is odd.}$

**Remark 3.2**: Equality in (13) holds only if  $n_1 = n_2 = k+1$  and  $2 \le k \le 5$ .

**Theorem 3.2**: For  $k \neq 1$ , there is no edge-maximal graph in  $\mathcal{C}(n,k)$  having a distinguished edge.

100

**Proof**: Let  $G \in \mathcal{C}(n,k)$ ,  $k \neq 1$ , be a graph containing a distinguished edge. To prove the theorem it is sufficient to show that G has less than  $\varepsilon(n,k)$  edges, where  $\varepsilon(n,k)$  is given by (11).

As in the proof of Theorem 3.1, we select a k-edge cut  $E_1$  containing a distinguished edge  $e_1$  such that  $G - E_1$  contains a component,  $G_1$ say, having no distinguished edge. Let  $G_2$  be the other component of  $G - E_1$  and let  $n_1 = |V(G_1)|$ , i = 1, 2. Since  $E_1$  contains a distinguished edge, then  $n_1 \ge k+1$ , and of course  $n_2 \ge k+1$ .

Since  $G_1$  has no distinguished edge of G Lemma 2.3 implies  $\varepsilon(G_1) \leq \varepsilon(n_1,k)$ . Consequently, if  $\varepsilon(G_2) \leq \varepsilon(n_2,k)$  then

$$\varepsilon(G) = \varepsilon(G_1) + \varepsilon(G_2) + k$$
  

$$\leq \varepsilon(n_1, k) + \varepsilon(n_2, k) + k,$$
  

$$\leq \varepsilon(n_1 + n_2, k) = \varepsilon(n, k). \qquad (Lemma 2.5)$$

Since  $k \ge 2$  and  $n_1 \ge k+1$  for i=1,2, it follows from Remark 3.1 that the above holds with equality only when  $n_1 = n_2 = k+1$  and  $2 \le k \le 4$ . Since  $\varepsilon(k+1,k) = \left\lfloor \frac{1}{8}(2k+1)^2 \right\rfloor = \frac{1}{2}k(k+1)$ , we conclude that  $G_1 = G_2 = K_{k+1}$ . Now since there are k edges between  $G_1$  and  $G_2$  we can assume, without loss of generality, that  $G_1$  contains two vertices,  $v_1$  and  $v_2$  say, joined to vertices in  $G_2$ . But then, since  $\kappa'(G_1) = k$ , there are at least k+1 edge-disjoint  $(v_1, v_2)$  - paths in  $G_1$ , contradicting the fact that G is k-critical. Hence  $\varepsilon(G) < \varepsilon(n,k)$  when  $\varepsilon(G_2) \le \varepsilon(n_2,k)$ .

Assume then that  $\varepsilon(G_2) > \varepsilon(n_2,k)$ . Then  $G_2$  contains at least one distinguished edge. As in the proof of Theorem 3.1, let T be the

largest segment of G such that  $T \cap V(G_1) = \phi$  and  $\varepsilon(G[T]) \le \varepsilon(|T|,k)$ . Let  $T' = V(G_2) \setminus T$ . Now  $T' \ne \phi$  since  $\varepsilon(G_2) > \varepsilon(n_2,k)$ . As in the proof of Theorem 3.1, if  $\varepsilon(G[T'] \le \alpha(|T'|,k)$ , then

$$\varepsilon(G) \leq \varepsilon(G_1) + \varepsilon(G[T]) + \varepsilon(G[T']) + 2k$$

$$\leq a(n_{k},k) + a(|T|,k) + a(|T'|,k) + 2k.$$

If  $|T'| \ge 2$ , then by Lemma 2.5

$$\varepsilon(G) \leq a(n,k) + a(|T| + |T'|,k) + k$$

Since  $n_1 \ge k+1$  and  $|T| + |T'| \ge k+3$  we have by Lemma 2.5 and Remark 3.1

$$a(n_1,k) + a(|T| + |T'|,k) + k$$

$$< a(n_{+} + |T| + |T'|,k),$$

and hence

$$\varepsilon(G) < a(n_1 + |T| + |T'|,k) = \varepsilon(n,k),$$

as required.

Consider now |T'| = 1. Then a(|T'|,k) = 0, and hence

 $\varepsilon(G) \leq a(n_1,k) + a(|T|,k) + 2k.$ 

If  $n_1 \ge 3k$  or  $|T| \ge 3k$ , then by Lemma 2.5 and Remark 3.1,

$$a(n_1,k) + a(|T|,k) + 2k < a(n_1 + |T| + 1,k).$$

Hence

$$\epsilon(G) < a(n_1 + |T| + 1,k) = \epsilon(n,k),$$

as required. So suppose that  $n_1 < 3k$  and |T| < 3k. Then, by Lemma 2.6

$$a(n_{1},k) + a(|T|,k) + 2k \le a(n_{1} + |T| + 1,k)$$

with equality holding only if  $n_1 = |T| = k+1$  and  $2 \le k \le 5$  (by Remark 3.2). As before, since  $a(k+1,k) = \frac{1}{2}k(k+1)$ ,  $G_1 = G[T] = K_{k+1}$ . Let v be the single vertex T'. Since  $d_G(v) \ge k+1$  ( $\ge 3$ ), v is joined to at least two vertices, say x and y in the same segment  $G_1$  or T of G. But then, there are at least k+1 edge-disjoint (x,y) - paths in G, contradicting the criticallity of G. Consequently

$$\varepsilon(G) < a(n_1 + |T| + 1, k)$$
$$= a(n, k) = \varepsilon(n, k)$$

as required. This completes the proof of the theorem for the case  $\varepsilon(G[T']) \leq \alpha(|T'|,k)$ .

Now suppose that  $\varepsilon(G[T']) > a(|T'|,k)$ . Then  $|T'| \ge k+1$ . As in the proof of Theorem 3.1, we partition T' into sets  $B_1, B_2, \ldots, B_t$  such that each subgraph  $G[B_i]$  contains no distinguished edge of G and hence

$$\varepsilon(G[B_i] \le a(|B_i|,k))$$

If  $|B_i| \ge k+1$  for some i, say i=1, then

 $\varepsilon(G) \leq \varepsilon(G_1) + \varepsilon(G[T]) + \varepsilon(G[B_1]) + 2k + \sum_{i=2}^{t} \varepsilon(G[B_i]) + (t-1)k$ 

$$\leq a(n_{1},k) + a(|T|,k) + a(|B_{1}|,k) + 2k + \sum_{i=2}^{t} a(|B_{i}|,k) + (t-1)k$$

$$\leq a(n_1,k) + a(|T|+|B_1|,k) + k + \sum_{i=2}^{t} a(|B_i|,k) + (t-1)k$$

(Lemma 2.5).

Since  $n_1 \ge k+1$  and  $|T| + |B_1| \ge 2(k+1)$ , then by Lemma 2.5 and Remark 3.1 we have

$$a(n_1,k) + a(|T| + |B_1|,k) + k < a(n_1 + |T| + |B_1|,k),$$

and hence since  $n_1 + |T| + |B_1| \ge 3(k+1)$ ,

$$\epsilon(G) < a(n_1, + |T| + |B_1|, k) + \sum_{i=2}^{t} a(|B_i|, k) + k(t-1)$$

$$\leq a(n,k) = \varepsilon(n,k),$$
 (Lemma 2.5)

as required.

Finally, when  $|B_i| \le k$  for every i,  $1 \le i \le t$ , the desired contradiction is obtained by applying the arguement use for the corresponding case in the proof of Theorem 3.1. This completes the proof of the Theorem.

Before proceeding to the characterisation, we need to describe the following graphs. Let H'(n,r) be a r-regular (r-semi-regular) graph on n vertices for even (odd) n.

Let n and k be two integers with  $2 \le k+1 \le n < 3k$ . We define

$$n_1 = \lfloor \frac{n+k+1}{2} \rfloor$$

and, for n > k+1, we construct the graph G' such that

$$G' = \overline{K}_{n-n_1} \vee H'(n_1, k + n_1 - n).$$

If  $n_1(k + n_1 - n)$  is even, then we define  $G_1^* = G'$ ; otherwise we define

$$\begin{split} G_1^* &= G' - \{e\}, \text{ where } e = xy \in E(G') \text{ with } x \in V(\overline{K}_{n-n_1}) \text{ and } y \in V(H'(n_1,k)) \\ &+ n_1 - n) \text{ such that } d_{H'}(y) = k-n+n_1+1. \quad \text{For } n=k+1, \text{ we take } G_1^* = H(n,k). \\ &\text{ Now we define } \hat{n}_1 = \left\lceil \frac{n+k+1}{2} \right\rceil, \quad \text{if } n = 3k - 2(2i-1), \text{ where } i = 1, 2, \dots, \frac{1}{4}(2k+1). \end{split}$$
We construct the graph  $G_2^*$  as follows :

$$G_2^* = H(\hat{n}_1, k)$$
, if  $n = \hat{n}_1$ , and  $G_2^* = K_{n-\hat{n}_1} \vee H'(\hat{n}_1, k+\hat{n}_1-n)$ , if  $n > \hat{n}_1$ .

Now for

$$n = 2i - k + 1;$$
  $k + 1 \le i \le 2k - 1$ 

or

n = 2i - k;  $k + 2 \le i \le i \le 2k - 1;$  i is odd,

let

$$\underline{n}_1 = \left\lfloor \frac{n+k-1}{2} \right\rfloor$$

Construct graph G such that

 $G = \overline{K}_{n-\underline{n}_{1}} \vee H'(\underline{n}_{1}, k + \underline{n}_{1} - n)$ 

We define  $G_3^* = G$  if  $\underline{n}_1(k - n + \underline{n}_1)$  is even and  $G_3^* = G - \{e\}$  if  $\underline{n}_1(k - n + \underline{n}_1)$  is odd where  $e = xy \in E(G)$ , such that  $x \in V(\overline{K}_{n-\underline{n}_1})$  and  $y \in V(H')$  with  $d_G(y) = k + 1$ .

Observe that for  $k + 1 \le n < 3k$  the graphs  $G_1^*$ ,  $G_2^*$  and  $G_3^*$ , are in the class  $\mathcal{A}(n,k)$  and have  $\left\lfloor \frac{1}{8}(n+k)^2 \right\rfloor$  edges and hence are edge-maximal in the class  $\mathcal{A}(n,k)$ .

Our next result provides us a characterisation of the edge-maximal graph in C(n,k).

**Theorem 3.3 :** G is an edge-maximal graph in  $\mathcal{C}(n,k)$ ,  $k \neq 1$  if and only if

(i) 
$$G = K_{k,n-k}$$
, if  $n \ge 3k$   
(ii)  $G = K_{k,n-k}$  or  $G = G_1^*$ , if  $n = 3k-1$   
(iii)  $G = K_{k,n-k}$  or  $G = G_1^*$  or  $G = G_2^*$ , if  $n = 3k-2$   
(iv)  $G = G_1^*$  or  $G = G_2^*$ , if  $n = 3k - 2$  (2i-1),  
 $i = 2, 3, ..., \frac{1}{4}(2k+1)$   
(v)  $G = G_1^*$  or  $G = G_3^*$ , if  $n = 2i - k + 1$ ,  $k + 1 \le i \le 2k - 1$   
or  $n = 2i - k$ , i is odd,  $k + 2 \le i \le 2k - 1$   
(vi)  $G = G_1^*$ , otherwise.

**Proof** : Let G be an edge-maximal graph in  $\mathcal{C}(n,k)$ ;  $k \neq 1$ . Then by Theorem 3.2, G has no distinguished edge and so  $G \in \mathcal{A}(n,k)$ . In other words, every edge e of G is incident to at least one vertex of degree k. Since  $G \in \mathcal{C}(n,k)$ , then  $\delta(G) = k$ . Now let  $G \in \mathcal{A}(n,k)$ . We denote by  $X = \{v \in V(G) : d_G(v) = k\}$  and  $\overline{X} = V(G) \setminus X$ . Let  $|X| = n_1$  and  $n_2 = |\overline{X}| = n - n_1$ . Now since every vertex v of G in  $\overline{X}$  has  $d_G(v) > k$ , and G  $\in \mathcal{A}(n,k)$ ,  $n_1 \geq k+1$ .

If  $n_1 \leq n-k$  then it is obvious that

$$\varepsilon(G) \leq n_1 k = g(n_1) .$$

Clearly

$$\max_{n} \{g(n_{1})\} = g(n-k) = k(n-k)$$

Now  $n_1 = n-k$ , if and only if  $n_2 = n - n_1 = k$  and so  $K_{k,n-k}$  is the only

edge-maximal graph for the case  $n_1 = n-k$ .

If  $n_1 \ge n-k$ , then  $n_2 = n - n_1 \le k$  simple counting gives :

$$\varepsilon(G) \leq \begin{cases} n_1 (n-n_1) + \left\lceil \frac{1}{2}n_1(k-n-n_1) \right\rceil , \text{ if } n_1(k-n+n_1) \text{ even} \\ \\ n_1(n-n_1) - 1 + \left\lceil \frac{1}{2}n_1(k-n+n_1) \right\rceil, \text{ if } n_1(k-n+n_1) \text{ odd} \end{cases}$$

That is

$$\varepsilon(G) \leq g(n_1) = n_1(n - n_1) + \left\lfloor \frac{1}{2}n_1(k - n + n_1) \right\rfloor.$$
  
Since for  $n \geq 3k$ ,  $g(n_1)$  is decreasing in  $n_1 \geq n - k$  we have

$$\max \{g(n_1)\} = g(n - k)$$

$$\prod_{i=1}^{n} = (n - k) k + \left\lfloor \frac{1}{2}(n - k)(0) \right\rfloor$$

$$= (n - k)k.$$

Again for this case  $K_{k,n-k}$  achieves this bound. Now for n < 3k, it is easy to verify that  $g(n_1)$  is increasing in  $n_1$  for  $n_1 \le \lfloor \frac{n+k-1}{2} \rfloor$ and decreasing in  $n_1$  for  $n_1 \ge \lfloor \frac{n+k+1}{2} \rfloor$ . Thus the maximum of  $g(n_1)$  is attained when

$$n_1 = \lfloor \frac{n+k+1}{2} \rfloor$$

Some simple algebra gives

$$g\left(\left\lfloor \frac{n+k+1}{2}\right\rfloor\right) = \left\lfloor \frac{1}{8}(n+k)^2\right\rfloor$$

Observe that the graph G achieving this bound is  $G_1^*$ . Now when  $n_1 = \lfloor \frac{n+k-1}{2} \rfloor$ , we have

$$g\left(\left\lfloor \frac{n+k-1}{2}\right\rfloor\right) = \left\lfloor \frac{1}{8}(n+k)^2 \right\rfloor$$

only if n = 2i - k + 1 for  $k + 1 \le i \le 2k - 1$  or n = 2i - k for  $k + 2 \le i \le 2k - 1$  and i is odd. So for this case the graph  $G_3^*$  achieves this bound. It can be shown that

$$g\left(\left\lceil \frac{n+k+1}{2}\right\rceil\right) = \left\lfloor \frac{1}{8} \left(n+k\right)^2 \right\rfloor$$

only if n = 3k - 4i+2, i = 1,2,..., $\frac{1}{4}(2k+1)$ . Clearly for this case  $G_2^*$  achieves this bound. Now since

$$(n-3k)^2 = (n+k)^2 - 8k(n-k) \ge 0$$

and k(n-k) is integer, we have

$$\left\lfloor \frac{1}{8}(n+k)^2 \right\rfloor \ge k(n-k)$$

for n < 3k, with equality holding only if n = 3k-1 or n = 3k-2. So for this case if G is an edge maximal then G could be  $K_{k,n-k}$  as well.

This completes the proof of the Theorem 3.3.

Note that for k = 1, every tree on n vertices belongs to C(n,k), having n-1 edges. Thus it is obvious that every tree on n vertices is an edge-maximal graph in C(n,1).

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